

BIGGER, BADDER BUGS

ABSTRACT. In this paper we motivate the ‘principles of trust’, chance-credence principles that are strictly stronger than the New Principle yet strictly weaker than the Principal Principle, and argue, by proving some limitative results, that the principles of trust conflict with Humean Supervenience.

1. INTRODUCTION

Humean Supervenience is the speculative, albeit appealing, thesis that the nomic supervenes on the categorical.¹ This paper asks whether Humean Supervenience is compatible with there being a tight enough connection between chance and rational credence, and offers new reasons for thinking not.

Past work is instructive.² There is, on the one hand, some familiar bad news for Humeans: Humean Supervenience is incompatible with the Principal Principle. In fact, Humean Supervenience is incompatible with the weakening of the Principal Principle one gets by a restriction to initial chance and rational initial credence. If Ch is the initial chance function, Cr is the class of rational initial credence functions, p is a proposition, and $\langle Ch(p) = x \rangle$ is the proposition that the initial chance of p equals x , then we have the following, a principle that asserts that rational initial credence *reflects* initial chance:³

Reflection. $\forall \pi \in Cr : \pi(p \mid \langle Ch(p) = x \rangle) = x$.

As the so-called ‘big, bad bug’ shows, Humean Supervenience and Reflection are not both true if chance has the features that science takes it to have. (See §3 for more.)

There is, on the other hand, some familiar good news for Humeans: Humean Supervenience is compatible with the New Principle.⁴ Restricting the New Principle to initial chance and rational initial credence gives us the following, a principle that asserts that rational initial credence *new-reflects* initial chance:

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¹Like Briggs (2009a), we take Humean Supervenience to be necessary and *a priori* if true, distinguishing it from the thesis that the nomic supervenes on the distribution of the categorical properties which are intrinsic to point-sized regions or objects, which may be, as Vranas (2002) and Lewis (1994) argue, contingent and/or *a posteriori*. Although, for reasons discussed in footnote 13, that assumption may not be necessary.

²The literature discussing Humean Supervenience and chance-credence principles is vast; see e.g. (Arntzenius and Hall, 2003), (Bigelow et al., 1993), (Briggs, 2009a,b), (Hall, 1994, 2004), (Halpin, 1994, 1998), (Hicks, 2017), (Ismael, 2008), (Levinstein, 2023), (Lewis, 1980, 1994), (Pettigrew, 2012, 2015, 2016), (Schaffer, 2003), (Thau, 1994), (Vranas, 2002), and (Ward, 2005).

³Let Ch_t be the chance function that holds at time t , and let q be any proposition. The Principal Principle is the following: $\forall \pi \in Cr : \pi(p \mid q \wedge \langle Ch_t(p) = x \rangle) = x$, if q is admissible w.r.t. $\langle Ch_t(p) = x \rangle$. See (Lewis, 1980).

⁴Let Ch_t be the chance function that holds at time t , and let q be any proposition. Then we have the New Principle: $\forall \pi \in Cr : \pi(p \mid q \wedge \langle Ch_t(p \mid q) = x \rangle) = Ch_t(p \mid q \wedge \langle Ch_t(p \mid q) = x \rangle)$. See (Hall, 1994), (Lewis, 1994), and Thau (1994).

22 **New Reflection.** $\forall \pi \in Cr : \pi(p \mid \langle Ch(p) = x \rangle) = Ch(p \mid \langle Ch(p) = x \rangle)$.

23 Chance having the features science takes it to have does not force a choice between
24 Humean Supervenience and New Reflection. (See §4 for more.)

25 Past work leaves much undecided, however. New Reflection does not draw a
26 tight enough connection between chance and credence. And a case can be made
27 that Reflection is stronger than need be: that the connection between chance and
28 rational credence can be tight enough, even if Reflection fails. An investigation of
29 intermediate chance-credence principles, strictly stronger than New Reflection and
30 strictly weaker than Reflection, is thus prompted.

31 This paper focuses primarily on three such: collectively, the *principles of trust*.⁵
32 The first asserts that rational initial credence *simply trusts* initial chance:

33 **Simple Trust.** $\forall \pi \in Cr : \pi(p \mid \langle Ch(p) \geq x \rangle) \geq x$.⁶

34 The second strengthens Simple Trust by ensuring that a rational initial credence
35 function updated on some information simply trusts initial chance updated on the
36 same information, thus asserting that rational initial credence *resiliently trusts* initial
37 chance:

38 **Resilient Trust.** $\forall \pi \in Cr : \pi(p \mid q \wedge \langle Ch(p \mid q) \geq x \rangle) \geq x$.⁷

39 The third, strictly stronger than the previous two, strengthens Simple Trust by
40 extending it to the expectation of all random variables. If χ is a random variable,
41 $\mathbb{E}_\pi(\chi)$ is the expectation of χ derived from some rational initial credence function
42 π , and $\mathbb{E}_{Ch}(\chi)$ is the expectation of χ derived from Ch , then we have the following,
43 a principle that asserts that rational initial credence *totally trusts* initial chance:

44 **Total Trust.** $\forall \pi \in Cr : \mathbb{E}_\pi(\chi \mid \langle \mathbb{E}_{Ch}(\chi) \geq x \rangle) \geq x$.⁸

45 Simple Trust and Resilient Trust may be easier to grok, but Total Trust is the principle
46 of greater interest, the principle demarcating the jointier epistemic joint. Some
47 properties that chance ought to have — some properties that chance must have,
48 we claim, if the connection between chance and rational credence is tight enough —
49 are had by chance only if Total Trust holds. (See §6 for more.)

50 Reflection is substantially stronger than Total Trust, as recent work on higher-
51 order evidence underscores. A case can be made that rational initial credence,
52 though not reflecting itself, totally trusts itself.⁹ Hoping that Humean Supervenience
53 will prove compatible with Total Trust, despite being incompatible with Reflection,
54 is thus — prior to a proper investigation of the matter — not unreasonable.

55 But the new news is bad news for Humeans. The compatibility of Humean
56 Supervenience and Total Trust is doubtful. In fact, in light of the limitative results
57 proved below, it is doubtful that any rational initial credence function totally trusts
58 initial chance if Humean Supervenience holds. One of the bigger, badder bugs
59 below concerns Simple Trust. We develop an argument that no rational initial
60 credence function simply trusts initial chance if Humean Supervenience holds. But
61 the assumptions of that argument are stronger than are the assumptions needed for

⁵For discussion of intermediate chance-credence principles, including the principles of trust, see e.g. (Dorst, 2019, 2020), (Dorst et al., 2021), (Elga, 2013), (Levinstein, 2023). Also see (Schervish, 1989).

⁶Equivalently, using upper bounds: $\forall \pi \in Cr : \pi(p \mid \langle Ch(p) \leq x \rangle) \leq x$.

⁷Equivalently, using upper bounds: $\forall \pi \in Cr : \pi(p \mid q \wedge \langle Ch(p \mid q) \leq x \rangle) \leq x$.

⁸Equivalently, using upper bounds: $\forall \pi \in Cr : \mathbb{E}_\pi(\chi \mid \langle \mathbb{E}_{Ch}(\chi) \leq x \rangle) \leq x$.

⁹This case is made in (Dorst, 2019, 2020) and (Dorst et al., 2021).

62 the other bigger, badder bug: the argument that no rational initial credence function
63 totally trusts chance if Humean Supervenience holds.

64 2. INVENTORY OF FORMAL TOOLS

65 Let us begin with an inventory of the formal tools invoked below.

66 There is, to begin with, a set of possible worlds, W , assumed (for convenience)
67 to be finite, and a set of propositions, identified with the powerset of W .¹⁰

68 There is also a set of random variables. A *random variable* χ is a function that maps
69 each possible world w to some real number, $\chi(w)$, the value of χ at w . One special
70 set of random variables is the set of indicator variables, the random variables whose
71 only possible values are 0 and 1. The set of indicator variables is, in a certain sense,
72 interchangeable with the set of propositions: for each indicator variable χ , there is
73 a unique proposition that contains world w just if $\chi(w) = 1$; for each proposition p ,
74 there is a unique indicator variable that maps world w to 1 just if w is an element
75 of p .

76 There is the aforementioned set of rational initial credence functions, Cr . Every
77 credence function maps each proposition to some real number on the unit interval,
78 and we assume that every rational initial credence function is a regular probability
79 function: a function that satisfies the probability axioms and gives nonzero credence
80 to every nonempty proposition.¹¹ Rational credence evolves: a rational agent's
81 present credence is arrived at by conditioning their rational initial credence function
82 on the information they have gathered heretofore. But to keep things simple, we
83 set non-initial credence aside, letting 'credence' hereafter denote initial credence.

84 There is also the *chance assignment*, a function that maps each world w to the
85 initial chance function that holds at w , namely, Ch_w . We assume that every possible
86 initial chance function is a probabilistic credence function. Chance evolves: the
87 present chances are arrived at by conditioning the initial chance function on the
88 history of the world heretofore. But to keep things simple, we set non-initial chance
89 aside, letting 'chance' hereafter denote initial chance.

90 Uncertainty about chance is uncertainty about chance *de dicto*. If an agent is
91 uncertain whether the chance of p equals x , they are not uncertain, for any world
92 w , about whether $Ch_w(p) = x$. What they are uncertain about is whether $Ch(p) = w$:
93 whether the chance of p , whatever it is, equals x . Claims about chance are thus,
94 unless otherwise noted, claims about chance *de dicto*. The proposition that the (*de*
95 *dicto*) chance of p equals x , $\langle Ch(p) = x \rangle$, is a set that includes world w just if $Ch_w(p) = x$;
96 the proposition that the (*de dicto*) chance of p is at least x , $\langle Ch(p) \geq x \rangle$, is a set that
97 includes world w just if $Ch_w(p) \geq x$.

98 Random variables are not bearers of chance; only propositions are. But ran-
99 dom variables have (*de dicto*) chance-expectations, and our space of propositions
100 includes propositions concerning the chance-expectations of random variables.
101 The *chance-expectation* of χ , $\mathbb{E}_{Ch}(\chi)$, is a Ch -weighted average of the possible values
102 of χ , $\sum_{v \in W} Ch(v)\chi(v)$. The proposition that the chance-expectation of χ equals x ,
103 $\langle \mathbb{E}_{Ch}(p) = x \rangle$, is a set that includes world w just if $\sum_{v \in W} Ch_w(v)\chi(v) = x$; the proposi-
104 tion that the chance-expectation of χ is at least x , $\langle \mathbb{E}_{Ch}(p) \geq x \rangle$, is a set that includes
105 world w just if $\sum_{v \in W} Ch_w(v)\chi(v) \geq x$.

¹⁰To ease the exposition, we ignore the distinction between a world and its singleton.

¹¹Assuming that every rational initial credence function is regular simplifies many of the arguments below. But the assumption is not essential.

106

3. THE BIG, BAD BUG

107 With the inventory of formal tools behind us, let us rehearse the big, bad bug:
 108 an argument that the conjunction of Humean Supervenience and Reflection is
 109 inconsistent with scientific practice.

110 Humean Supervenience is a constraint on the chance assignment. Possible worlds
 111 can be partitioned by their Humean mosaics.¹² A cell of the partition is a *mosaic*.
 112 A chance assignment verifies Humean Supervenience just if it maps any pair of
 113 worlds in the same mosaic to the same chance function.¹³

114 Reflection is another constraint on the chance assignment. The chance assignment
 115 verifies Reflection only if some rational credence function reflects the chances it
 116 engenders. A chance assignment is *immodest* just if it verifies the following, a
 117 principle that asserts that each possible chance function gives itself chance one:

118 **Immodesty.** For any worlds v and w , if $Ch_v \neq Ch_w$, then $Ch_v(w) = 0$.

119 And, if we ignore degenerate chance assignments (as we will, hereafter), Reflection
 120 implies Immodesty: a regular probability functions reflects the chances engendered
 121 by a non-degenerate chance assignment only if the chance assignment is
 122 immodest.¹⁴

123 There are chance assignments that verify both Humean Supervenience and
 124 Immodesty, but there is a third constraint. An adequate chance assignment must
 125 accord with scientific practice. It is not easy to say what it takes to accord with
 126 scientific practice, but a necessary condition is ready to hand. Consider *the best-*
 127 *system function*: a function that maps each mosaic to the theory or theories that
 128 best systematize the mosaic, as judged by the method of theory choice implicit in
 129 science. Any theory that could be among the outputs of the best-system function
 130 determines a chance function over the space of possible worlds. A chance function
 131 *systematizes* a mosaic just if it is determined by all of the theories to which the
 132 best-system function maps the mosaic. To accord with scientific practice, a chance
 133 assignment must verify:

134 **Possible Systematization.** Every chance function is compossible
 135 with every mosaic it systematizes.

136 Verifying Possible Systematization is easy if Humean Supervenience fails, since
 137 different chance functions then can hold at worlds in the same mosaic. But if Humean
 138 Supervenience holds, then a chance function is compossible with a mosaic only if it
 139 is necessitated by the mosaic. Humean Supervenience and Possible Systematization
 140 thus together imply:

141 **Necessary Systematization.** Every chance function is necessitated
 142 by every mosaic it systematizes.

¹²Or anyway one must assume to take Humean Supervenience seriously.

¹³Here we rely on the assumption Humean Supervenience is necessary if true. For a defense of the assumption, see (Briggs, 2009a, 443-44). But insofar as we are interested in Resilient Trust or Total Trust, the assumption is not essential. If Humean Supervenience is contingent, then we can focus on the following claim entailed by Resilient Trust: every rational initial credence function updated on Humean Supervenience simply trusts chance update on Humean Supervenience.

¹⁴Reflection is a norm of local chance reflection. There is also a norm of global chance reflection: $\forall \pi \in Cr : \pi(p \mid \langle Ch = Ch_w \rangle) = Ch_w(p)$. The global norm straightforwardly implies Immodesty; see Fact 3.1 of (Dorst, 2020, 616). And although, strictly speaking, the local and global norms are not equivalent, the difference between them can be ignored. For, as Gallow (2023) proves, they come apart only in the degenerate case in which the chance assignment is 'half-cyclic'.

143 A chance function is *system-modest* just if it assigns positive chance to a mosaic
 144 systematized by a distinct chance function. If some mosaic is systematized by a
 145 system-modest chance function, then Immodesty and Necessary Systematization
 146 are not both true. Possible Systematization, Humean Supervenience, and Immodesty
 147 together imply:

148 **Immodest Systematization.** No mosaic is systematized by a system-
 149 modest chance function.

150 And therein lies the problem, for Immodest Systematization is false. There is room
 151 for disagreement about when a chance function systematizes a mosaic. The method
 152 of theory choice implicit in science is not entirely transparent to us. But nor is
 153 it entirely opaque. We know enough about it to know that some mosaics are
 154 systematized by system-modest chance functions.

155 There are realistic ways of illustrating the failure of Immodest Systematization.
 156 Lewis (1994, 482) appeals to radioactive decay, noting that a mosaic systematized
 157 by a chance function that encodes one half-life for a given radioactive particle gives
 158 positive chance to mosaics systematized by distinct possible chance functions that
 159 encode distinct half-lives for the same radioactive particle. But partly to make the
 160 problem clearer and partly to set the stage for the limitative results below, we will
 161 appeal to, as we call them, ‘flip models’.

162 Each flip model is associated with some natural number, n . The mosaic of a
 163 world in an n -flip model is a binary sequence of length n , envisaged, picturesquely,
 164 as the outcomes of the flips of some quantum coin: *HTHHTH* We assume that
 165 every binary sequence of length n is the mosaic of some world in the n -flip model;
 166 we assume — identifying worlds and mosaics and thereby hardcoding the truth
 167 of Humean Supervenience — that no binary sequence of length n is the mosaic of
 168 more than one world in the n -flip model; and we assume each world w has some
 169 precise chance function, Ch_w .¹⁵ We thus can refer to an n -flip model as a pair $\langle W, \mathcal{P} \rangle$,
 170 where W is the set of binary sequences of length n , and \mathcal{P} is a function from W to
 171 probability functions over W , i.e., $\mathcal{P} : W \rightarrow \Delta(W)$, $w \mapsto Ch_w$.

172 We call a chance function IID when it treats the coin flips as independent and
 173 identically distributed. Formally, if H_j is the proposition that the j^{th} flip lands heads,
 174 then:

175 **IID.** Chance function Ch is IID just if, for any j and k , $j < k \leq n$:

176 (1) $Ch(H_j \wedge H_k) = Ch(H_j)Ch(H_k)$, and

177 (2) $Ch(H_j) = Ch(H_k)$.

178 One expects the chances associated with coin flips to be distributed binomially,
 179 and it is the IID chance functions that deliver binomial distributions. Let $IID(x)$
 180 be the IID chance function *centered* on x , the chance function that deems each flip
 181 independent and accords each flip chance x of landing heads; and let $\langle Ch = IID(x) \rangle$
 182 be the proposition that holds at world w just if $Ch_w = IID(x)$. If w is a world in the
 183 n -flip model at which $\langle Ch = IID(x) \rangle$ holds, and v is a world in the n -flip model at
 184 which k of the n flips land heads, then $Ch_w(v) = x^k(1-x)^{(n-k)}$; hence if $\langle \#H = k \rangle$ is the
 185 proposition that exactly k of the n flips land heads, $Ch_w(\langle \#H = k \rangle) = \binom{n}{k} x^k (1-x)^{(n-k)}$.

186 Some venerable approaches to chance entail that every world in a flip model is
 187 systematized by an IID chance function. For example, according to frequentism,

¹⁵For some w , Ch_w may be deterministic, i.e., it may specify result of each flip with probability 1.

188 whenever exactly k of the n flips at world w land heads, $Ch_w = IID(k/n)$. Frequentism
 189 is not obvious, however. Consider the following, from the 20-flip model:

190 $w_i : HHHHHHHHHHTTTTTTTTTT$

191 It may be that the best-system function maps w_i to the deterministic theory that
 192 a flip lands if and only if it is among the first ten flips, in which case the chance
 193 function that systematizes w_i does not treat the coin flips as identically distributed.

194 But we know that many worlds in flip models are systematized by IID chance
 195 functions — IID chance processes are ubiquitous in science, the norm from which
 196 exceptions deviate. We know that the following, from the 20-flip model, is system-
 197 atized by $IID(1)$:

198 $w_j : HHHHHHHHHHHHHHHHHHHHH$

199 We know that the following, from the 20-flip model, is systematized by $IID(0)$:

200 $w_k : TTTTTTTTTTTTTTTTTT$

201 And we know that many of the worlds wherein exactly half of the flips land heads
 202 are systematized by $IID(1/2)$, the following being a good candidate:

203 $w_l : HTHTHTTTTHHTTTHTTHH$

204 Arguably, we know something stronger. The great virtue of focusing on flip
 205 models is that it allows to state precise claims about what science requires of the
 206 chance assignment, and a case can be made that we know the following, a principle
 207 that asserts that $IID(x)$ systematizes some world in the n -flip model whenever x
 208 is the actual proportion of heads to flips at some world in the model:

209 **Proportional Systematization.** For any m and n , $0 \leq m \leq n$, there is
 210 some world in the n -flip model systematized by $IID(m/n)$.

211 Proportional Systematization is plausible and interesting, and it will play an im-
 212 portant role in one of the bigger, badder bugs to come.

213 But if our aim is only to bring out the falsity of Immodest Systematization,
 214 nothing so strong is needed. Indeed, the following suffices:

215 **Nontrivial Systematization.** In some n -flip model, some world is
 216 systematized by $IID(x)$, $0 < x < 1$, and some world is systematized
 217 by some chance function distinct from $IID(x)$.

218 Nontrivial Systematization is an extremely weak claim about what science requires
 219 of a chance assignment, yet it is inconsistent with Immodest Systematization. If some
 220 world in the n -flip model is systematized by $IID(x)$, and some world is systematized
 221 by a chance function distinct from $IID(x)$, then every world systematized by $IID(x)$
 222 is systematized by a system-modest chance function, since $IID(x)$ gives positive
 223 chance to every world in the n -flip model.

224 Taking a step back, we can see the structure of the gauntlet facing Humeans.
 225 The big, bad bug has three parts. There is a scientific part, a purported claim
 226 about what science requires of the chance assignment. There is an epistemological
 227 part, the claim that the connection between chance and rational credence is tight
 228 enough only if Reflection holds. And there is the mathematical part, a proof that
 229 Humean Supervenience is inconsistent with Reflection, given the purported claim
 230 about what science requires of the chance assignment. Humeans wax poetic about
 231 the epistemological virtues of their metaphysics, the optimific balance of strength,
 232 simplicity, and fit that chance and laws as they envisage them achieve. But the big,
 233 bad bug is an impossibility result, and waxing poetic is not adequate response to an

234 impossibility result. What Humeans need is a tenability result: a *proof* that Humean
 235 Supervenience is consistent with some not-too-loose connection between chance
 236 and rational credence, given some not-too-weak claim about what science requires
 237 of the chance assignment.

238 4. NEW REFLECTION

239 The gauntlet facing Humeans would be less formidable if New Reflection drew
 240 a tight enough connection between chance and rational credence. But it doesn't.

241 Indeed, New Reflection bears on the connection between chance and rational
 242 credence only indirectly. What it directly bears on is the connection between
 243 rational credence and, as we will call it, 'informed chance'. For each possible chance
 244 function Ch_w , there is the proposition that Ch_w holds, $\langle Ch = Ch_w \rangle$, and the informed
 245 chance function at world w , Ch_w^+ , is $Ch_w(- \mid \langle Ch = Ch_w \rangle)$, the chance function at
 246 w conditioned on $\langle Ch = Ch_w \rangle$. Our space of propositions includes propositions
 247 concerning the (*de dicto*) informed chances of propositions. The proposition that
 248 the informed chance of p equals x , $\langle Ch^+(p) = x \rangle$, is a set that includes world w just if
 249 $Ch_w^+(p) = x$; the proposition that the informed chance of p is at least x , $\langle Ch^+(p) \geq x \rangle$,
 250 is a set that includes world w just if $Ch_w^+(p) \geq x$.

251 New Reflection is equivalent to the following, a principle that asserts that rational
 252 credence *reflects* informed chance:

253 **Informed Reflection.** $\forall \pi \in Cr : \pi(p \mid \langle Ch^+(p) = x \rangle) = x$.

254 The connection New Reflection draws is thus just as tight as the connection Reflec-
 255 tion draws, but whereas Reflection connects rational credence and chance, New
 256 Reflection connects rational credence and informed chance.

257 If chance is immodest, then chance and informed chance coincide: $Ch_w = Ch_w^+$
 258 for each world w . But if Humean Supervenience holds, then chance is modest,¹⁶
 259 and if chance is modest, then chance and informed chance can come apart.

260 A frequentist, 2-flip model provides a simple illustration. There are four worlds,
 261 HH , HT , TH , and TT . If frequentism holds at each, then $Ch_{HH} = IID(1)$, $Ch_{HT} =$
 262 $IID(1/2) = Ch_{TH}$, and $Ch_{TT} = IID(0)$. But the chance of both flips landing heads is
 263 $1/4$ only if exactly one flip land heads. So chance and informed chance come apart:
 264 for example, $Ch_{HT}(HH) = 1/4 < Ch_{HT}^+(HH) = 0$.

265 The connection New Reflection draws between rational credence and informed
 266 chance induces an indirect connection between chance and rational credence. But
 267 the induced connection is not tight enough if chance and informed chance can come
 268 apart, as we can see by considering anti-expertise.

269 Say that credence function π treats *de dicto* probability function P as an *anti-expert*
 270 with respect to some proposition-value pair, (p, x) , just if $\pi(p \mid \langle P(p) \geq x \rangle) < x$ and
 271 $\pi(p \mid \langle P(p) < x \rangle) \geq x$; and say that P is *free of anti-expertise* just if no rational credence
 272 function treats P as an anti-expert with respect to any proposition-value pair. While
 273 Reflection entails that chance is free of anti-expertise,¹⁷ New Reflection does not. In
 274 fact, it is consistent with New Reflection that chance is rife with anti-expertise.

275 Chance is, as Lewis says, a guide to life:

¹⁶Humean Supervenience, Possible Systematization, and the negation of Immodest Systematization together entail the negation of Immodesty.

¹⁷Chance is free of anti-expertise if and only if Simple Trust holds, and Reflection entails Simple Trust.

276 It is reasonable to let one's choices be guided in part by one's firm
 277 opinions about objective chances or, when firm opinions are lacking,
 278 by one's degrees of belief about chance. . . . The greater chance you
 279 think the ticket has of winning, the greater should be your degree
 280 of belief that it will win; and the greater is your degree of belief that
 281 it will win, the more, *ceteris paribus*, it should be worth to you and
 282 the more you should be disposed to choose it over other desirable
 283 things. (1980, 287-88)

284 But because it is consistent with New Reflection that chance is rife with anti-
 285 expertise, it is consistent with New Reflection that chance is an anti-guide to life. It
 286 is consistent with New Reflection that rational agents often take truth and chance
 287 to be anti-correlated, regarding as evidence against p information that increases
 288 what they think the chance of p is. It is thus consistent with New Reflection that
 289 rational agents often prefer a lesser chance to a greater chance of getting the things
 290 they desire. And that, we think, is absurd. Chance is not an anti-guide to life; and
 291 from that we conclude that every tight enough chance-credence principle entails
 292 that chance is free of anti-expertise.

293 A two-world model provides an illustration. Suppose that w and v each accords
 294 the other more chance than it accords itself: $Ch_w(v) = Ch_v(w) = 0.9$, and $Ch_w(w) =$
 295 $Ch_v(v) = 0.1$. The agent prefers w to v . The agent divides their credence equally
 296 between the two worlds and new-reflects chance: $\pi(w) = \pi(v) = 0.5$, and for any
 297 p , $\pi(p \mid \langle \langle Ch(p) = x \rangle \rangle) = Ch(p \mid \langle \langle Ch(p) = x \rangle \rangle)$. The agent then regards chance as an
 298 anti-expert: the agent thinks that evidence that the chance of w is low is evidence
 299 that w is true, and thus prefers a lesser chance of getting what they prefer, a lesser
 300 chance of w , to a greater chance. See Figure 1 for a depiction of this scenario.

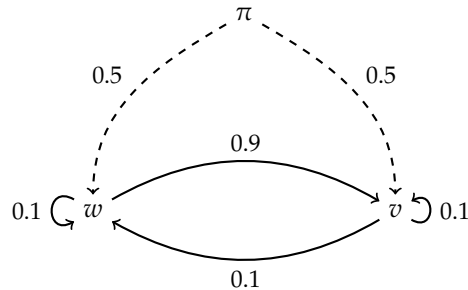


FIGURE 1. π assigns w and v probability .5. $Ch_w(v) = Ch_v(w) = .9$,
 and $Ch_w(w) = Ch_v(v) = .1$. π new reflects Ch .

301 Reflection is tight enough — Reflection entails that chance is free of anti-expertise.
 302 But Reflection implies Immodesty, and as the big, bad bug shows, Humean Supervenience
 303 is incompatible with any chance-credence principle that entails Immodesty.
 304 The principles of trust thus prove their interest; for all of them entail that chance is
 305 free of anti-expertise, and none of them imply Immodesty.¹⁸

¹⁸Here is a two-world model that verifies Simple Trust and falsifies Immodesty: $Ch_w(w) = Ch_v(v) = 0.8$; $Ch_w(v) = Ch_v(w) = 0.2$.

306

5. SIMPLE TRUST

307 Simple Trust, the weakest of the principles of trust, is equivalent to the claim that
 308 chance is free of anti-expertise. So if every tight enough chance-credence principle
 309 entails that chance is free of anti-expertise, Simple Trust holds.

310 Simple Trust also can be motivated by appeal to accuracy. Say that credence
 311 function π treats *de dicto* probability function P as *expectedly inaccurate* just if, for
 312 some acceptable way of measuring accuracy, π expects itself to be more accurate
 313 than P ; and say that P is *free of expected inaccuracy* just if no rational credence function
 314 treats P as expectedly inaccurate. Chance ought to be free of expected inaccuracy.
 315 The indicator function at world w specifies the value of each indicator variable
 316 at w , thus, given the aforementioned interchangeability of indicator variables and
 317 propositions, specifying the truth-value of each proposition at w . Chance is highly
 318 inaccurate at world w just if the divergence between Ch_w and the indicator function
 319 at world w is great, and while proponents and opponents of Humean Supervenience
 320 disagree about the prevalence of worlds at which chance is highly inaccurate, all
 321 sides agree that no rational (initial) credence function gives high credence to worlds
 322 at which chance is highly inaccurate.

323 Chance is free of expected inaccuracy only if Simple Trust holds, however. In fact,
 324 the implication goes both ways. As Levinstein (2023) shows, if the received view is
 325 correct about what the acceptable ways of measuring are — if the acceptable ways
 326 of measuring accuracy are the additive, strictly proper, truth-directed measures that
 327 satisfy certain continuity and limit assumptions — then Simple Trust is equivalent
 328 to the claim that chance is free of expected inaccuracy.¹⁹

329

6. TOTAL TRUST

330 It is doubtful that Simple Trust is itself tight enough, however, for two reasons.

331 The first concerns accuracy. The accuracy argument for Simple Trust, when gen-
 332 eralized, becomes an argument for Total Trust. The specification function at world
 333 w generalizes the indicator function at world w , specifying the value of all random
 334 variables at w . A probability function P induces an estimate function, \mathbb{E}_P , which
 335 maps each random variable χ to some real number, $\mathbb{E}_P = \sum_w P(w)\chi(w)$, and just as
 336 divergence is distance between probability and indication, estimate inaccuracy —
 337 the generalization of inaccuracy to all random variables — is divergence between
 338 estimate and specification. The estimate inaccuracy for a set of random variables
 339 of probability function P at world w is a measure of how far apart \mathbb{E}_P is from the
 340 specification function for those variables at w .²⁰

341 Say that credence function π treats *de dicto* probability function P as *expectedly*
 342 *estimate inaccurate* just if, for some acceptable way of measuring estimate inaccuracy,
 343 π expects itself to be more estimate accurate than P for some random variable; and
 344 say that P is *free of expected estimate inaccuracy* just if no rational credence function
 345 expects itself to be more expectedly estimate accurate than P for any random
 346 variable. Chance ought to be free of expected estimate inaccuracy, for the same
 347 reasons that chance ought to be free of expected inaccuracy. But, as Dorst et al.
 348 (2021) show, generalizing the result proved in Levinstein (2023), if the received view
 349 is correct about what the acceptable ways of measuring estimate inaccuracy are —

¹⁹For the precise conditions required on measures of accuracy, see (Levinstein, 2023).

²⁰For technical details, see (Dorst et al., 2021) and Campbell-Moore (MS).

350 if the acceptable measures of estimate inaccuracy are strictly proper, truth-directed
 351 measures that satisfy certain continuity and limit assumptions — then Total Trust
 352 is equivalent to the claim that chance is free of expected estimate inaccuracy.²¹

353 The second reason concerns choice. If chance is a guide to life, then deferring a
 354 choice to chance — letting chance choose on one’s behalf, as it were; giving chance
 355 power of attorney — ought always to be rational. But deferring a choice to chance
 356 is always rational only if Total Trust holds. In fact, the implication holds both ways.
 357 As Dorst et al. (2021) show, Total Trust is equivalent to the claim that deferring a
 358 choice to chance is always rational.

359 *Choice technicalities:* A choice is a set of pairwise exclusive options, $O = \{o_1, \dots, o_n\}$.
 360 Each option is a random variable, a function that maps each world to some real
 361 number which represents how desirable the agent finds the option at the world.
 362 The expected value of option o , relative to credence function π , $V(\pi, o)$, equals
 363 $\sum_W \pi(w)o(w)$.

364 Deferring a choice among O to chance is a strategy: the chance-expected value of
 365 option o at world v is $\sum_w Ch_v(w)o(w)$, and deferring a choice among O to chance is
 366 a function that maps each world v to some option that maximizes chance-expected
 367 value at v . If $s(w)$ is the value at w of the option to which world w is mapped by
 368 the strategy of deferring a choice to chance, then the expected value of deferring a
 369 choice among O to chance, relative to credence function π , is $\sum_W \pi(w)s(w)$.

370 Credence function π *permits* deferring a choice among O to P just if, for each o in
 371 O , $V(\pi, o) \leq V(\pi, s)$. It is *rational* to defer a choice among O to P just if every rational
 372 credence function permits deferring a choice among O to P . And it is *always* rational
 373 to defer a choice to P just if, for any O , it is rational to defer a choice among O to P .
 374 *End of choice technicalities.*

375 It is doubtful that the connection between chance and rational credence is tight
 376 enough if it is not always rational to defer a choice to chance. Deferring a choice to
 377 chance is playing the chances, selecting an option that maximizes chance-expected
 378 value, and if chance is a guide to life, then it should always be rational to play the
 379 chances. But if it is always rational to defer a choice to chance, then Total Trust
 380 holds: the claim that every rational credence function totally trusts some *de dicto*
 381 probability function P is equivalent to the claim that it is always rational to defer a
 382 choice to P .²²

383 It is an interesting question whether Total Trust is itself tight enough. One
 384 worry stems from expectation-matching.²³ Another worry stems from stochastic

²¹For the precise statement and proof of this result, see (Dorst et al., 2021)

²²Dorst et al. (2021) offer an example to help illustrate the difference between Simple Trust and Total Trust. Suppose that there are three worlds, w , v , and u . Suppose that there are two options, $o_0(w) = o_0(v) = o_0(u) = 0$, $o_1(w) = 29$, $o_1(v) = -3$, and $o_1(u) = -13$. And consider the following chance assignment: $Ch_w(w) = 0.45$, $Ch_w(v) = 0.10$, and $Ch_w(u) = 0.45$; $Ch_v(w) = 0.15$; $Ch_v(v) = 0.70$, and $Ch_v(u) = 0.15$; and $Ch_u(w) = 0.30$, $Ch_u(v) = 0.10$, and $Ch_u(u) = 0.60$. At each of the three worlds, the chance-expected value of o_1 exceeds zero, and hence exceeds the chance-expected value of o_0 . But some probabilistic credence functions that simply trusts (and indeed resiliently trusts) this chance assignment nevertheless strictly prefer o_0 to o_1 . One example is $\pi(w) = 0.17$, $\pi(v) = 0.56$, and $\pi(u) = 0.27$.

²³Matching one’s credences to one’s expectation of the chances is a central part of science and an ubiquitous part of daily life. It is thus insist that a chance-credence principle entail Chance Expectation: $\forall \pi \in Cr : \pi(p) = \sum_W \pi(w)Ch_w(p)$. Reflection entails Chance Expectation, but Total Trust does not, as the following two-world model illustrates: $\pi(v) = \pi(w) = 0.5$; $Ch_v(v) = 0.9$; $Ch_v(w) = 0.1$; $Ch_w(w) = 0.8$; and $Ch_w(v) = 0.2$; cf. (Dorst et al., 2021, n. 18).

385 dominance.²⁴ But what is relevant for our argumentative purposes is the necessity
 386 claim, not the sufficiency claim, and the case that every tight enough chance-credence
 387 principle entails Total Trust is strong.

388 7. A BIGGER, BADDER BUG

389 Our first limitative result concerns Simple Trust. Consider the following, a prin-
 390 ciple that asserts that every proposition is compossible with every possible propo-
 391 sition that sets a positive lower bound on its chance:

392 **Threshold Compossibility.** For every value $x > 0$, if $\langle Ch(p) \geq x \rangle$ is
 393 possible, then $p \wedge \langle Ch(p) \geq x \rangle$ is possible.

394 Simple Trust entails Threshold Compossibility. In fact, no regular probability func-
 395 tion simply trusts chance if Threshold Compossibility fails.²⁵ And as flip models
 396 make clear, the conjunction of Humean Supervenience and Threshold Compossibil-
 397 ity is incompatible with plausible claims about what science requires of the chance
 398 assignment. For example, as we prove in this section, in any n -flip model, $n > 4$,
 399 Threshold Compossibility is incompatible with Proportional IID.

400 The proof proceeds by cases. Let a k -heads world be a world at which k flips land
 401 heads, and consider the following, a principle that asserts that IID(k/n) holds at
 402 some world w in an n -flip model only if w is a k -heads world:

403 **Matching.** For any world w in an n -flip model, if $Ch_w = IID(x/n)$,
 404 then $w \in \langle \#H = x \rangle$.

405 If Proportional IID holds, and Matching fails, then the chance that some world
 406 accords itself is exceeded by the chance accorded to it by some other world. To see
 407 this, take an arbitrary counterinstance to Matching: suppose that $Ch_w = IID(k/n)$,
 408 and suppose that w is a j -heads world, $j \neq k$. Since Proportional IID holds, there is
 409 some world v in the n -flip model at which IID(j/n) holds. For any z , $0 \leq z \leq n$, the
 410 chance of w at a world at which IID(z/n) holds equals $(z/n)^j(1 - (z/n))^{n-j}$, which
 411 takes its unique maximum at $z = j$. The chance of w at v thus exceeds the chance
 412 of w at w , and Threshold Compossibility therefore fails. The proposition that the
 413 chance of w is at least as high as the chance of w at v is, although possible, not
 414 compossible with w .

415 Threshold Compossibility also fails, however, in any n -flip model, $n > 4$, if
 416 Proportional IID and Matching hold, as we see clearly in the 6-flip model. Let
 417 $\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle$ be the proposition that the coin lands heads either exactly two

²⁴The proposition that the value of option o exceeds x ($o \geq x$), is a set that includes world w just if $o(w) \geq x$. The proposition that option o_i chance-wise stochastically dominates option o_j , $\langle o_i > o_j \rangle$, is a set that includes world w just if (a) for every x , $Ch_w(\langle o_i \geq x \rangle) \geq Ch_w(\langle o_j \geq x \rangle)$, and (b) for some x , $Ch_w(\langle o_i \geq x \rangle) > Ch_w(\langle o_j \geq x \rangle)$. Reasoning by chance-wise stochastic dominance is ubiquitous and intuitive. It is thus natural to insist that a chance-credence principle entail Chance-wise Stochastic Dominance: $\forall \pi \in Cr : \text{if } \pi(\langle o_i > o_j \rangle) > 0$, then $\sum_W \pi(w \mid \langle o_i > o_j \rangle) o_i(w) \geq \sum_W \pi(w \mid \langle o_i > o_j \rangle) o_j(w)$. Reflection entails that Chance-wise Stochastic Dominance, but Total Trust does not, as the following four-world model illustrates: $\pi(u) = \pi(v) = \pi(w) = \pi(x) = 1/4$; $\pi = Ch_u$; $Ch_u(u) = 2/9$, $Ch_u(v) = 1/3$, $Ch_u(w) = 2/9$, and $Ch_u(x) = 2/9$; $Ch_v(u) = 2/11$, $Ch_v(v) = 3/11$, $Ch_v(w) = 4/11$, and $Ch_v(x) = 2/11$; $Ch_x(u) = 2/13$, $Ch_x(v) = 3/13$, $Ch_x(w) = 4/13$, and $Ch_x(x) = 4/13$; $o_i(u) = 1$, $o_i(v) = 2$, $o_i(w) = 0$, and $o_i(x) = 4$; and $o_j(u) = 4$, $o_j(v) = 0$, $o_j(w) = 1$, and $o_j(x) = 1$. Although π totally trusts chance, $\sum_W \pi(w \mid \langle o_i > o_j \rangle) o_i(w) = 1.5 < \sum_W \pi(w \mid \langle o_i > o_j \rangle) o_j(w) = 2$.

²⁵If π is a rational credence function, and $\langle Ch(p) \geq x \rangle$ is possible, then $\pi(p \mid \langle Ch(p) \geq x \rangle)$ is defined. If $\pi(p \mid \langle Ch(p) \geq x \rangle)$ is defined, and $p \wedge \langle Ch(p) \geq x \rangle$ is impossible, then $\pi(p \mid \langle Ch(p) \geq x \rangle) = 0 < x$.

418 or exactly four times; let w_2 be a 2-heads world at which $IID(2/6)$ holds; let w_3 be a 3-
 419 heads world at which $IID(3/6)$ holds; and let w_4 be a 4-heads world at which $IID(4/6)$
 420 holds. Because of the bell-shape of the binomial curve, $Ch_{w_3}(\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle)$,
 421 the sum of the fairly high chance w_3 accords to 2-heads worlds and the fairly high
 422 chance w_3 accords to 4-heads worlds, exceeds both $Ch_{w_2}(\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle)$, the
 423 sum of the high chance w_2 accords to 2-heads worlds and the low chance w_2 accords
 424 to 4-heads worlds, and $Ch_{w_4}(\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle)$, the sum of the low chance that w_4
 425 accords to 2-heads worlds and the high chance that w_4 accords to 4-heads worlds.

$$Ch_{w_2}(\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle) = \binom{6}{2} \left(\frac{2}{6}\right)^2 \left(\frac{4}{6}\right)^4 + \binom{6}{2} \left(\frac{2}{6}\right)^4 \left(\frac{4}{6}\right)^2 \approx 0.41$$

$$Ch_{w_3}(\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle) = \binom{6}{2} \left(\frac{3}{6}\right)^2 \left(\frac{3}{6}\right)^4 + \binom{6}{2} \left(\frac{3}{6}\right)^4 \left(\frac{3}{6}\right)^2 \approx 0.47$$

$$Ch_{w_4}(\langle \#H = 2 \rangle \vee \langle \#H = 4 \rangle) = \binom{6}{2} \left(\frac{4}{6}\right)^2 \left(\frac{2}{6}\right)^4 + \binom{6}{2} \left(\frac{4}{6}\right)^4 \left(\frac{2}{6}\right)^2 \approx 0.41$$

426 For a visual depiction, see Figure 2.

427 We thus can produce a counterexample to Threshold Compossibility by taking
 428 any nonempty subset of $\langle Ch = IID(2/6) \rangle \vee \langle Ch = IID(4/6) \rangle$, which includes exactly
 429 as many elements of $\langle Ch = IID(2/6) \rangle$ as $\langle Ch = IID(4/6) \rangle$. One example is the
 430 disjunction of w_2 and w_4 :

$$Ch_{w_2}(w_2 \vee w_4) \approx 0.027$$

$$Ch_{w_3}(w_2 \vee w_4) \approx 0.031$$

$$Ch_{w_4}(w_2 \vee w_4) \approx 0.027$$

431 The calculations above pertain only to the 6-flip model. But similarly reasoning
 432 shows that in any n -flip model, $n > 4$, Threshold Compossibility fails if Proportional
 433 IID and Matching both hold.²⁶

434 Proportional IID enjoys considerable plausibility. If it is possible that a quantum
 435 coin flipped n times lands heads exactly m times, then it seems possible that each
 436 flip of a quantum coin flipped n times be independent and have chance m/n of
 437 landing heads. A Humean who denies Proportional IID thus denies the possibility
 438 of something that seems possible. Of course, Humeans are committed to denying
 439 the possibility of things that seem possible already. It seems possible that an
 440 indeterministic quantum coin lands heads on each of its n flips. But there is only
 441 one n -heads world in the n -flip model. So if a Humean thinks that the n -heads
 442 world in the n -flip model is deterministic, a world in which it is nomically necessary

²⁶For each m , let w_m be a m -heads world in the n -flip model at which $IID(m/n)$ holds. If $n > 4$ is even, then $w_{(n-2)/2} \vee w_{(n+2)/2}$ is not compossible with the claim that the chance of $w_{(n-2)/2} \vee w_{(n+2)/2}$ is at least x , where x is the chance of $w_{(n-2)/2} \vee w_{(n+2)/2}$ at $w_{n/2}$. If $n > 4$ is odd, then $w_{(n-3)/2} \vee w_{(n+1)/2}$ is not compossible with the claim that the chance of $w_{(n-3)/2} \vee w_{(n+1)/2}$ is at least x , where x is the chance of $w_{(n-3)/2} \vee w_{(n+1)/2}$ at $w_{(n-1)/2}$.

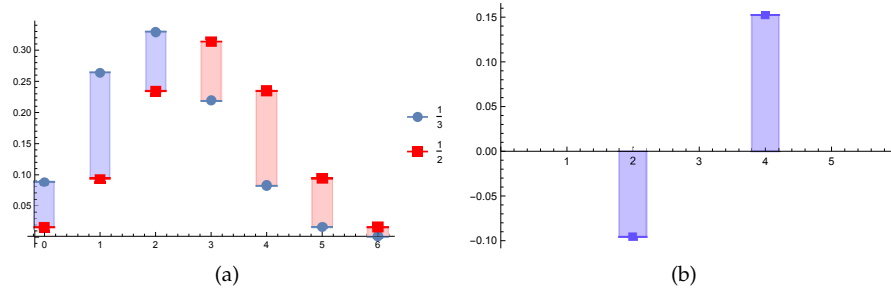


FIGURE 2. Figure 2(a) displays the probabilities assigned to 0, 1, 2, 3, 4, 5, 6 occurrences of heads for IID(1/2) and IID(1/3). Figure 2(b) isolates the difference assigned to two occurrences and six occurrences of heads. Although IID(1/2) assigns lower probability to there being exactly two occurrences of heads than IID(1/3) does, it assigns significantly higher probability to there being exactly four occurrences of heads.

443 that every flip lands heads, then the Humean must deny that it is possible that
 444 an indeterministic quantum coin land heads on each of its n -flips. But denying
 445 Proportional IID is not just denying the possibility of something that seems possible.
 446 It is one thing to set limits on how far apart the underlying chances and frequencies
 447 can be. It is another thing to set limits on how close together they can be. The
 448 chances that feature in our best scientific theories often are arrived at by fitting a
 449 curve to the actual frequencies.

450 And the full strength of Proportional IID is not needed to render Threshold
 451 Compossibility and Humean Supervenience incompatible. Say that x is a possible
 452 IID center in an n -flip model just if $IID(x)$ holds at some world in the n -flip model.
 453 The thrust of the point then can be put, vaguely but helpfully, as follows: *Threshold*
 454 *Compossibility fails in an n -flip model whenever the possible IID centers are sufficiently*
 455 *clustered*. Proportional IID entails that the possible IID centers are sufficiently
 456 clustered, but weakenings do likewise. For example: if there are three possible IID
 457 centers inclusively between 8/20 and 12/20 in the 20-flip model, then Threshold
 458 Compossibility fails.

459 Reconciling Simple Trust and Humean Supervenience is harder than reconciling
 460 Threshold Compossibility and Humean Supervenience — Threshold Compossibility
 461 does not entail Simple Trust. But appreciating the challenge of reconciling
 462 Threshold Compossibility and Humean Supervenience helps us see how formidable
 463 the gauntlet facing Humeans is. Science requires that there be many possible
 464 IID centers, and apparently weak claims about the diversity and distribution of
 465 possible IID chance in flip models renders Threshold Compossibility false.

466 8. ANOTHER BIGGER, BADDER BUG

467 The next limitative result concerns Total Trust. Consider the following, a principle
 468 that asserts that there are at least two nontrivial possible IID centers in big enough
 469 flip models.

470 **Nontrivial Diversity.** If n is big enough, then for some x and y ,
 471 $0 < x < y < 1$, $IID(x)$ and $IID(y)$ each hold at some world or other
 472 in the n -flip model.

473 There is a claim about the extent of IID chance: a claim, clarified and made precise
 474 below, about the proportion of worlds in flip models at which IID chance functions
 475 hold. The claim is weak — it is very plausible that its truth is part of what science
 476 requires of the chance assignment. And as we prove (in Appendix A), Nontrivial
 477 Diversity and Total Trust are not both true, if this weak claim about the extent of
 478 IID chance holds.

479 The tension Total Trust engenders in flip models between the extent of IID chance
 480 and the diversity of possible IID centers is easy to see if we consider a very strong
 481 claim about the extent of IID chance.

482 Call w and v *mirrored* in an n -flip model just in case the sequence of heads and
 483 tails in w and v is exactly switched. That is, H_j (heads on the j^{th} flip) holds at w just
 484 in case T_j holds at v . For example, in a five flip model, the world $HHTTH$ and the
 485 world $TTHHT$ are mirrored. The following constraint requires a symmetry between
 486 mirrored worlds when one has an IID chance function.

487 **Symmetry.** An n -flip model is *symmetric* just if, for all $w \in W$, if
 488 $Ch_w = IID(x)$, and v mirrors w , then $Ch_v = IID(1 - x)$.

489 Let $\#w$ be the number of occurrences of heads at w . I.e., $\#w = k$ just in case
 490 $w \in \langle \#H = k \rangle$. We then have the following result:

491 **Initial Triviality.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model totally trusted by some
 492 π , all members of \mathcal{P} are IID, and $\langle W, \mathcal{P} \rangle$ is symmetric, then if $0 <$
 493 $\#w < n$, $Ch_w = IID(1/2)$.

494 So, for example, if Total Trust holds, all of the possible chance functions in the
 495 1000-flip model are IID, and the 1000-flip model is symmetric, then Nontrivial
 496 Diversity fails; for $IID(1/2)$, then, holds at every world in the 1000-flip model,
 497 except perhaps the 0-heads and the 1000-heads world.²⁷

498 Here is a sketch of the proof:

499 *Proof.* (Sketch) The proof appeals to a background fact (theorem A.1
 500 in Appendix A): If $\langle W, \mathcal{P} \rangle$ is an n -flip frame, then some regular
 501 probability function π totally trusts $\langle W, \mathcal{P} \rangle$ if and only if the members
 502 of \mathcal{P} totally trust one another.

503 Suppose each element of \mathcal{P} is IID, and suppose that $\langle W, \mathcal{P} \rangle$ is
 504 symmetric. We show that if the elements of \mathcal{P} resiliently trust one
 505 another, then $Ch_w = Ch_v$ for all $Ch_w, Ch_v \in \mathcal{P}$ unless there are either
 506 0 or n occurrences of heads at w or v .

507 Let E be the proposition that there are either $n - 1$ or n total
 508 occurrences of heads and H^n be the proposition that all flips are
 509 heads. By Symmetry and the fact that all chance functions are IID,
 510 all worlds with the same number of occurrences of heads have the
 511 same chance function. Let Ch_j refer to the chance function at all

²⁷The idea for this result depends on the fact that, in a binomial distribution, the probability of all flips coming up heads decreases very rapidly for $IID(x)$ as x decreases. Suppose then, that Ch_w is $IID(x)$ for some low x . If Ch_w conditions on the fact that the chance of heads is actually *high*, it still won't assign high probability to all heads. That is, $Ch_w(\text{All heads} \mid Ch(H) \text{ is high})$ will still be too low.

512 worlds with j -occurrences of heads and let $Ch_j(H) = p_j$.²⁸ Finally, let
 513 $Ch_{n-1}(H^n | E) = x$.

We can then derive that:

$$(1) \quad \begin{aligned} Ch_1(H^n | E, \langle Ch(H^n | E) \geq x \rangle) &= Ch_1(H^n | E) \\ &= \frac{p_1^n}{np_1^{n-1}(1-p_1) + p_1^n} \end{aligned}$$

and

$$(2) \quad \begin{aligned} Ch_{n-1}(H^n | E, \langle Ch(H^n | E) \geq x \rangle) &= p_{n-1}(H^n | E) \\ &= \frac{p_{n-1}^n}{np_{n-1}^{n-1}(1-p_{n-1}) + p_{n-1}^n} \\ (3) \quad &= \frac{(1-p_1)^n}{n(1-p_1)^{n-1}p_1 + (1-p_1)^n} \\ &= x \end{aligned}$$

514 (Lines (1) and (2) follow from the fact that H is distributed according
 515 to a binomial distribution, and line (3) follows from Symmetry.)

516 If all functions in \mathcal{P} *totally* trust one another, then they *resiliently*
 517 *trust* one another. So, we check what is required to make line (1)
 518 greater than or equal to line (3). With some simple algebra, we find
 519 that this requires $p_1 \geq 1/2$ and $p_{n-1} \leq 1/2$. Given Resilient Trust, this
 520 entails that $p_1 = \dots = p_{n-1} = \frac{1}{2}$. \square

521 We prove a variant of this result in Appendix A (theorem A.12). Of course, even
 522 if the chance functions at many or most of the worlds in the n -flip model are IID,
 523 it is doubtful that every possible chance function in the n -flip model is IID. Initial
 524 Triviality thus puts little pressure, if any, on a Humean. But all that we need to
 525 render Nontrivial Diversity and Total Trust incompatible is a weak claim about the
 526 extent of IID chance: the claim, clarified and made precise immediately below, that
 527 the extent of IID chance in n -flip models does not decrease as n increases.

528 For simplicity, we consider only n -flip models where n is even, and we assume
 529 that there is at least one $(n/2)$ -heads world at which $IID(1/2)$ holds. We put these
 530 two ideas together with the following axiom:

531 **Fifty/Fifty.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model, then n is even, and at some
 532 $w \in \langle \#H = n/2 \rangle$, $Ch_w = IID(1/2)$.

533 It will now be useful to introduce some more definitions. For a given n -flip
 534 model, we say that a number m is in the **IID region** of n if there is some m -heads
 535 world at which an IID chance function holds. In notation, we write $IID(Ch_w)$ to
 536 mean Ch_w is IID, and we define $IID \text{ reg}(n) := \{m : \exists w \text{ s.t. } \#w = m \text{ and } IID(Ch_w)\}$.

537 We say that m is in the **even odds region** of n just if there is some m -heads
 538 world in the n -flip model at which $IID(1/2)$ holds. In notation, $EO\text{-region}(n) := \{m : \exists w \text{ s.t. } Ch_w = IID(1/2) \text{ and } \#w = m\}$. And we let $\ell(n)$ be the *smallest number* in the
 539 even odds region of n : $\ell(n) := \min_m m \in EO\text{-region}(n)$.
 540

²⁸For what we've said so far, some worlds with the same number of heads might still have (up to two) different IID chance functions. This slightly complicates the proof in tedious ways, so we omit details.

541 The next axiom codifies the earlier thought that IID chance functions are possible
 542 at worlds with a reasonable mixture of heads and tails. The specific assumption we
 543 need is:

544 **Sufficiency.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model, then (1) for all k such that
 545 $\frac{n}{4} \leq k \leq \frac{n}{2}$, k is in the IID region of n , and (2) if 0 is not in the even
 546 odds region of n , then $\ell(n) - 1$ is in the IID region of n .

547 The first part of this axiom ensures that an IID chance function holds at some k -heads
 548 world, if k is between $n/4$ and $n/2$. This seems very reasonable, especially in large
 549 models. There is, taking such a case, some 250,000-heads world in the 1,000,000-
 550 flip model without any discernible pattern beyond the fact that tails occurs three
 551 times as often as heads.²⁹ The second part ensures that there is some world with
 552 an IID chance function centered on something other than $1/2$ unless the model is
 553 completely trivial and assigns an IID chance function centered on $1/2$ even in the
 554 n -heads world.

555 The next assumption establishes a particular type of lower bound on the per-
 556 centage of worlds with IID chance functions.

Boundedness. There exists $d > 0$ and $N \in \mathbb{N}$ such that for all $n \geq N$,
 if $\langle W, \mathcal{P} \rangle$ is an n -flip model and m is in the IID region of n , then

$$\frac{|\{w : \#w = m \text{ and } IID(Ch_w)\}|}{|\{w : \#w = m\}|} \geq d$$

557 Here's the intuition. We let the Humean pick some number n that she counts as
 558 'big'. We also let her pick some really small lower bound. For concreteness, say big
 559 numbers are at least 100 and the lower bound is 1%. We give her a big n -flip model
 560 and ask her for which $m \leq n$ there is at least one m -heads world at which an IID
 561 chance function holds. This axiom then requires that at least one percent of the
 562 m -heads worlds have IID chance functions. She is free to make 'big' be as large as
 563 she likes, and she is free to make d be as small as she likes so long as it is bigger
 564 than 0.³⁰

565 This axiom is technical, but innocuous. Worlds at which IID chance functions
 566 hold are *disorganized*. There is not much to say about them beyond roughly what
 567 the frequency of heads to tails is. (If there were more to say, then there would be a
 568 nice law characterizing the pattern.) As n grows large, more and more worlds are
 569 disorganized — most sequences appear totally random. Think of a television screen
 570 with its mix of black and white pixels. There are a few arrangements of such pixels
 571 that result in discernible patterns, something you could relatively easily describe.
 572 But for the vast majority, the screen is just random noise. Denying Boundedness is
 573 akin thinking that discernible patterns are more common as size of the television
 574 screens increases, which is exactly the opposite of what seems clear. Discernible
 575 patterns are less common as the size increases.

576 The final axiom is required for technical reasons:

577 **Monotonicity.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model and Ch_w, Ch_v in \mathcal{P} are both
 578 IID with $Ch_w(H) < Ch_v(H)$, then $Ch_w(\langle Ch(H^n) \geq 2^{-n} \rangle) < Ch_v(\langle Ch(H^n) \geq$
 579 $2^{-n} \rangle)$.

²⁹When combined with Symmetry, Sufficiency guarantees there an IID chance function holds at some k -heads world, if k is between $n/2$ and $3n/4$.

³⁰We can actually weaken this axiom so that it only applies to m in the even odds region of n instead of in the IID region of n , but it strikes us as a bit less natural when stated that way.

580 If all worlds in the model are IID, then Monotonicity is redundant. In that case,
 581 $\langle Ch(H^n) \geq 2^{-n} \rangle = \langle Ch(H) \geq 1/2 \rangle$. This axiom rules out strange situations where
 582 many non-IID worlds with relatively few heads for some reason give fairly high
 583 probability to the claim that all flips land heads.

584 We can now state our most powerful result (see Appendix A for proof).

585 **Serious Triviality.** Let $\langle W_1, \mathcal{P}_1 \rangle, \langle W_2, \mathcal{P}_2 \rangle, \dots$ be a sequence of models
 586 with $|W_i| < |W_{i+1}|$. Assume each validates Sufficiency, Fifty/Fifty,
 587 Monotonicity, and Symmetry. Moreover, assume that Boundedness
 588 holds of the sequence. Then there exists an $N \in \mathbb{N}$ such that if $i \geq N$
 589 and some regular probability function totally trusts $\langle W_i, \mathcal{P}_i \rangle$, then
 590 for all $Ch_w \in \mathcal{P}_i$ such that $IID(Ch_w)$, we have $P_w = IID(1/2)$.

591 Serious Triviality tells us that the weak claim about the extent of IID chance — the
 592 conjunction of Sufficiency, Fifty/Fifty, Monotonicity, Symmetry, and Boundedness
 593 — implies that Total Trust and Nontrivial Diversity are not both true. If any rational
 594 credence function totally trusts chance and the weak claim about the extent of IID
 595 chance holds, then, for large n , every possible IID chance function in the n -flip
 596 model is centered on $1/2$, except possibly the 0-heads and n -heads worlds.

597 Science requires both that the extent of IID chance be considerable and that the
 598 diversity of possible IID centers be many. The case for Total Trust is strong. But as
 599 the proof of Serious Triviality reveals, no chance assignment that is totally trusted
 600 by a rational credence function provides both the extent of IID chance and the
 601 diversity of possible IID centers that science requires.

602 9. CONCLUSION

603 The big, bad bug shows that Humean Supervenience is inconsistent with Re-
 604 flection, given a hard-to-deny claim about what science requires of the chance
 605 assignment. A promising Humean response is to reject Reflection in favor of some
 606 principle that draws a looser but still tight enough connection between chance and
 607 credence. The connection that New Reflection draws is, we argue, not tight enough,
 608 so we are led to the principles of trust, intermediate principles, which are strictly
 609 weaker than Reflection yet strictly stronger than New Reflection. The suspicion that
 610 Humean Supervenience is not consistent with a tight enough connection between
 611 chance and credence would be greatly reduced with a tenability result: a proof
 612 that Humean Supervenience is consistent with some or all of the principles of
 613 trust, given some not-too-weak assumptions about what science requires of the
 614 chance assignment. But what we have instead are bigger, badder bugs: proofs that
 615 Humean Supervenience is inconsistent with principles of trust, given stronger but
 616 still hard-to-deny claims about what science requires of the chance assignment.

617 Our limitative results pertain to particularly simple flip models: finite, fixed flip
 618 models, wherein each world has the same number of flips. Some of our results extend
 619 to finite, variable flip models, wherein different worlds have different numbers of
 620 flips.³¹ But there is more work to do investigating both finite, variable flip models
 621 and infinite flip models.³²

³¹For example: the fact that Simple Trust and Proportional IID are not both true in a finite, fixed flip model implies that Resilient Trust and Proportional IID are not both true in a finite, variable flip model.

³²There is also work to do investigating flip models in which some worlds lack a precise chance function.

622 And there is work to do extending the argument beyond flip models. Realistic
 623 hypotheses about the world we find ourselves in are, in various ways, unlike a
 624 world exhausted by a sequence of coin flips. Even if a realistic hypothesis about our
 625 world could be encoded in a binary sequence, it is unlikely that our best scientific
 626 theories would treat each bit in the binary sequence as the outcome of some IID
 627 chance process. But the difference between the worlds in flip models and realistic
 628 hypotheses about the world we find ourselves in does not obviously provide solace
 629 to Humeans. Our experience suggests that reconciling Humean Supervenience and
 630 the principles of trust becomes harder, not easier, as the size and the complexity of
 631 the model increases.

632 The way forward is gradual and mathematically precise, proceeding from less
 633 to more realistic models. Our limitative results are just some of the very many
 634 out there — there is a continent to explore. There are many claims about what
 635 science requires of the chance assignment worth considering and many intermediate
 636 chance-credence principles besides the principles of trust. The continent is sure to
 637 contain stronger limitative results than the ones proved here. Whether the continent
 638 also contains philosophically interesting tenability results remains to be seen. Is
 639 there any proof that Humean Supervenience is consistent with some tight enough
 640 connection between chance and credence, given the truth of hard-to-deny claims
 641 about what science requires of the chance assignment?

642

A. APPENDIX

643 In the appendix, we prove a variant of the Initial Triviality result (Theorem A.12)
 644 and prove the Serious Triviality result (Theorem A.15).

645 **A.1. Notation and Terminology.** As before, we use $\langle W, \mathcal{P} \rangle$ to refer to a generic
 646 n -flip model. We will switch to using $P_w \in \mathcal{P}$ to refer to the chance function at a
 647 world (instead of Ch_w and P to refer to the (*de dicto*) chance function—whatever it
 648 is—instead of Ch partly for reasons of notational compactness and partly because
 649 the results hold generically for all such models even when P and P_w are interpreted
 650 differently.

651 As before, we will use loose talk and say that a function P_w is IID when it treats
 652 the flips in a sequence as IID. Even more loosely, we'll say a world w is IID just in
 653 case P_w is IID.

654 As in the main text, we will write $P_w = IID(x)$ to mean P_w is IID and assigns
 655 probability x to heads. It will also sometimes be convenient, when P_w is IID, to
 656 write $P_w(H) = x$ or $P_w(H) \geq x$. As in the main text, we will also write $IID(P_w)$ to
 657 mean that P_w is IID.

658 We'll say that $\langle W, \mathcal{P} \rangle$ **validates** Total/Simple Trust just in case all members of
 659 \mathcal{P} totally/simplely trust P . More explicitly, $\langle W, \mathcal{P} \rangle$ validates Simple Trust if for all w ,
 660 $P_w(p \mid \langle P(p) \geq x \rangle) \geq x$ for all x , and similarly for Total Trust.

661 As a reminder, we also have the following notation:

- 662 • $\#w$ refers to the number of heads at w .
- 663 • $\ell(n)$ is the smallest number k in an n -flip model obeying Fifty/Fifty such
 664 that for all w where $\#w = k$, w has an IID chance function centered on $1/2$.
- 665 • H^n refers to the proposition that all n flips in an n -flip model land heads.

666 We also remind the reader of the following principles for reference below (now
 667 with P and P_w replacing Ch and Ch_w).

668 **Symmetry.** An n -flip model is *symmetric* just if, for all $w \in W$, if
 669 $P_w = \text{IID}(x)$, and v mirrors w , then $P_v = \text{IID}(1 - x)$.

670 **Fifty/Fifty.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model, then n is even, and at some
 671 $w \in \langle \#H = n/2 \rangle$, $P_w = \text{IID}(1/2)$.

672 **Sufficiency.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model, then (1) for all k such that
 673 $\frac{n}{4} \leq k \leq \frac{n}{2}$, k is in the IID region of n , and (2) if 0 is not in the even
 674 odds region of n , then $\ell(n) - 1$ is in the IID region of n .

675 **Monotonicity.** If $\langle W, \mathcal{P} \rangle$ is an n -flip model and P_w, P_v in \mathcal{P} are both IID
 676 with $P_w(H) < P_v(H)$, then $P_w(\langle P(H^n) \geq 2^{-n} \rangle) < P_v(\langle P(H^n) \geq 2^{-n} \rangle)$.

Boundedness. There exists $d > 0$ and $N \in \mathbb{N}$ such that for all $n \geq N$,
 if $\langle W, \mathcal{P} \rangle$ is an n -flip model and m is in the IID region of n , then

$$\frac{|\{w : \#w = m \text{ and } \text{IID}(P_w)\}|}{|\{w : \#w = m\}|} \geq d$$

677 **A.2. Results.** Our main question concerns when a regular probability function
 678 π can totally trust chance. As it turns out, to answer that question, we just need
 679 to find out when the frame $\langle W, \mathcal{P} \rangle$ validates total trust, as the following theorem
 680 establishes.

681 **Theorem A.1** (Dorst et al.). *A regular probability function π totally trusts a frame $\langle W, \mathcal{P} \rangle$*
 682 *only if $\langle W, \mathcal{P} \rangle$ validates total trust.*

683 The proof is involved, so we omit it here and refer the interested reader to (Dorst
 684 et al., 2021, Theorem 4.1). As we'll see, our results below entail that the functions
 685 in \mathcal{P} can't even *simply* trust one another. We conjecture that no regular probability
 686 function can simply trust them.

687 For what follows, it's important to keep in mind that if $\text{IID}(P)$, then according to
 688 P, H follows a Bernoulli Distribution with parameter $P(H)$. In turn, if X is a random
 689 variable representing the total number of heads, then X is distributed according to
 690 a Binomial Distribution with parameter $P(H)$. If $P(H) = p$, the probability of any
 691 given world with $\#w = k$ is $p^k(1 - p)^{n-k}$. So, if $0 < p < 1$, then for all $w \in W$, $P(w) > 0$.

692 We now prove some basic facts about models that validate Simple Trust. (Dorst
 693 (2020) provides a more general result implying part (2) of the following proposition.)

694 **Proposition A.2.** *Suppose $\langle W, \mathcal{P} \rangle$ validates Simple Trust. Then*

- 695 (1) *If $\langle W, \mathcal{P} \rangle$ validates Fifty/Fifty, then for all $w \in W$, $P_w(w) > 0$, and*
 696 (2) *For all $w, v \in W$, $P_w(w) \geq P_v(w)$*

697 *Proof.* To prove (1): Let $P_h = \text{IID}(1/2)$ be in \mathcal{P} . (Existence is guaranteed by Fifty/Fifty).
 698 For all $w \in W$, it's clear $P_h(w) > 0$. Suppose $P_w(w) = 0$ for some $w \in W$. Then
 699 $w \in \langle P(w) \leq 0 \rangle$, so $P_h(w \mid \langle P(w) \leq 0 \rangle)$ is defined and > 0 . Contradiction.

700 To prove (2): Suppose $P_w(w) < P_v(w) = x$. Then $w \notin \langle P(w) \geq x \rangle$. So, $P_v(w \mid \langle P(w) \geq$
 701 $x \rangle) = 0 < x$. □

702 **Proposition A.3.** *Suppose $\langle W, \mathcal{P} \rangle$ validates Simple Trust. Let $P_w, P_v \in \mathcal{P}$ be IID with*
 703 *$\#w < \#v$. Then $P_w(H) \leq P_v(H)$.*

704 *Proof.* Let $P_w(H) = p_w$ and $P_v(H) = p_v$. Suppose $\#w < \#v$ but $p_w > p_v$. Recall that
 705 if X is the number of heads, then according to both P_v and P_w , X is distributed
 706 according to a Binomial Distribution with parameters p_v and p_w respectively. So, if
 707 $\frac{\#w}{n} \leq p_v < p_w$, then $P_v(w) > P_w(w)$, which entails $\langle W, \mathcal{P} \rangle$ violates Simple Trust, (by

708 part (2) of Proposition A.2). Likewise, if $p_v < \frac{\#w}{n} \leq p_w \leq \frac{\#v}{n}$, $P_w(v) > P_v(v)$. Finally,
 709 suppose $p_v < \frac{\#w}{n} \leq \frac{\#v}{n} \leq p_w$. In this case, $P_v(w) > P_v(v) \geq P_w(v) \geq P_w(w)$, again
 710 violating Simple Trust by Prop. A.2. \square

711 *Remark.* Proposition A.3 does not rule out the possibility of distinct IID chance
 712 functions at worlds w and v if $\#w = \#v$ in an n -flip model. Following the proof, we
 713 see there could be a maximum of two different IID chance functions for worlds
 714 with the same number of heads, namely, one on each side of $\#w/n$. (This adds a
 715 wrinkle elided over to the proof sketch of Initial Triviality in the main text, but it's
 716 one that can be easily accommodated.) As we'll now see, there is one important
 717 exception.

718 **Proposition A.4.** *Suppose $\langle W, \mathcal{P} \rangle$ validates Simple Trust and Fifty/Fifty. Then if $w \in W$
 719 is IID and $\#w = n/2$, $P_w = \text{IID}(1/2)$.*

720 *Proof.* By Fifty/Fifty, some world $h \in W$ is IID such that $\#h = n/2$ and $P_h = \text{IID}(1/2)$.
 721 So, if P_w is also IID and $\#w = n/2$, then $P_w(w) \leq P_h(w)$. Given Prop. A.2. $P_w(w) \geq$
 722 $P_h(w)$, so $P_w(H) = 1/2$. \square

723 *Remark.* Note that Proposition A.4 guarantees that for any n -flip model $\langle W, \mathcal{P} \rangle$
 724 validating Simple Trust and Fifty/Fifty, $\ell(n)$ is defined and $\leq \frac{n}{2}$. Further, we have
 725 $\ell(n) \geq 1$ by part 1 of Prop. A.2.

726 We can also put upper bounds on worlds with IID chance functions that have
 727 under $\ell(n)$ total heads.

728 **Fact A.5.** *Suppose $\langle W, \mathcal{P} \rangle$ is an n -flip model validating Simple Trust and Fifty/Fifty.
 729 Suppose P_w is IID for some world w with $\#w = \ell(n) - 1$. Then if $P_w \neq \text{IID}(1/2)$,
 730 $P_w(H) < \frac{\ell(n)}{n}$.*

731 *Proof.* Suppose $\langle W, \mathcal{P} \rangle$ validates Simple Trust and $\ell(n) \geq 1$. Let $\#w = \ell(n) - 1$, and
 732 let $P_w = \text{IID}(p)$. Suppose $\frac{\ell(n)}{n} \leq p$. Let $v \in W$ with $\#v = \ell(n)$ and $P_v = \text{IID}(1/2)$.
 733 By hypothesis, $p \neq 1/2$. By Prop. A.3, p must be $< 1/2$. But in that case, since
 734 $\frac{\ell(n)}{n} \leq p < 1/2$, $P_w(v) > P_v(v)$, contradicting Prop. A.2. \square

735 We know that $P_w(H) \leq P_v(H)$ if $\#w < \#v$ and both have IID chance functions
 736 by Proposition A.3. We also know, by Fact A.5 that if $\#w < \ell(n)$ and w is IID,
 737 $P_w(H) < \frac{\ell(n)}{n}$.

738 It will be useful below to consider a special IID probability function P^ℓ over
 739 W but *not* in \mathcal{P} such that $P^\ell(H) = \frac{\ell(n)}{n}$. The following lemma will serve to put an
 740 important constraint on P^ℓ . Namely, if $\langle W, \mathcal{P} \rangle$ validates Simple Trust and Fifty/Fifty,
 741 then $P^\ell(H \mid \langle H^n \geq 2^{-n} \rangle) \geq 2^{-n}$.

742 **Lemma A.6.** *Let $\langle W, \mathcal{P} \rangle$ be an n -flip frame validating Simple Trust with at least one
 743 IID function $P \in \mathcal{P}$ such that $P(H) \geq 1/2$. For any $x \in (0, 1)$, let $P^{(x)} = \text{IID}(x)$.³³ Let
 744 $f(x) = P^{(x)}(H^n \mid \langle P(H^n) \geq 2^{-n} \rangle)$. Then f is strictly increasing over $(0, 1)$.*

Proof. Let $V := \{w \in W \mid P_w(H^n) \geq 2^{-n}\}$. Note that the requirement that there be at
 least one IID chance function $P \in \mathcal{P}$ such that $P(H) \geq 1/2$ guarantees V is non-empty.

³³Note that $P^{(x)}$ is not necessarily in \mathcal{P} .

Let $V(k) := |\{w \in V : \#w = k\}|$. With f and $P^{(x)}$ defined as above, we then have

$$(4) \quad f(x) = \frac{x^n}{P^{(x)}(V)}$$

$$(5) \quad = \frac{x^n}{\sum_{k=0}^n V(k)x^k(1-x)^{n-k}}$$

745 f is clearly differentiable, so we just need to check that its derivative is positive.
 746 This is straightforward but tedious to do. \square

747 Our next goal is to put lower bounds on $\ell(n)$ for a given model (Lemma A.8). To
 748 do so must first prove Lemma A.7, which in turn appeals to the famous Inequality
 749 of Arithmetic and Geometric Means.

AM-GM Inequality. For any list of n non-negative reals x_1, \dots, x_n ,

$$\frac{1}{n} \sum_{i=1}^n x_i \geq \left(\prod_{i=1}^n x_i \right)^{\frac{1}{n}}$$

750 with equality iff $x_1 = x_2 = \dots = x_n$.

751 **Lemma A.7.** Suppose $n, k \in \mathbb{N}$ with $n > k$. Then $\frac{n}{2^{\frac{n+k}{n}}} \geq \frac{n-k}{2}$.

752 *Proof.* Simple algebra shows that the lemma holds if and only if for all $n \geq k + 2$,
 753 we have:

$$(6) \quad \frac{n}{n-k} \geq 2^{\frac{k}{n}}$$

To prove line (6), first consider a list of numbers x_1, \dots, x_n with:

$$x_i = \begin{cases} 2 & i \leq k \\ 1 & i > k \end{cases}$$

We have:

$$\frac{1}{n} \sum x_i = \frac{n+k}{n}$$

and

$$\left(\prod x_i \right)^{\frac{1}{n}} = 2^{\frac{k}{n}}$$

754 So, by the AM-GM Inequality, $2^{\frac{k}{n}} < \frac{n+k}{n}$.

755 To prove line (6) holds, we just need to determine when $\frac{n+k}{n} \leq \frac{n}{n-k}$, and it is easy
 756 to see this holds whenever $n > k$. \square

Let $\langle W, \mathcal{P} \rangle$ be an n -flip frame. Suppose $w \in W$ is a world with $\#w = \ell(n) - 1$ with
 IID chance function P_w . By Fact A.5, if $P_w(H) \neq 1/2$, $P_w(H) < \frac{\ell(n)}{n}$. Let P^ℓ be defined
 over W (but not necessarily in \mathcal{P}) such that $P^\ell = \text{IID}\left(\frac{\ell(n)}{n}\right)$. By Lemma A.6, we know

$$P_w(H^n \mid \langle P(H^n) \geq 2^{-n} \rangle) < P^\ell(H^n \mid \langle P(H^n) \geq 2^{-n} \rangle).$$

757 This will be important for the next lemma.

758 **Lemma A.8.** Suppose $\langle W, \mathcal{P} \rangle$ is an n -flip model validating Simple Trust, Fifty/Fifty, and
 759 Sufficiency, with $\ell(n) \geq 2$. Let $P^\ell(H) = \frac{\ell(n)}{n}$ be an IID probability function. Then (1)
 760 $P^\ell(\langle P(H^n) \geq 2^{-n} \rangle) > 0$ and (2) if $P^\ell(\langle P(H^n) \geq 2^{-n} \rangle) \geq 2^{-k}$ for $k \in \mathbb{N}$, then $\ell(n) \geq \frac{n-k}{2}$.

761 *Proof.* Part (1) follows trivially from the fact that $\ell(n) > 0$ and Fifty/Fifty.

762 We now establish part (2). Let $P_w \in \mathcal{P}$ be IID with $P_w(H) < \frac{1}{2}$ and $\#w = \ell(n) - 1 > 0$.
 763 Such a P_w is guaranteed to exist by Sufficiency. By Proposition A.2, $0 < P_w(H)$. Since
 764 P_w is also IID and $\langle W, \mathcal{P} \rangle$ validates Fifty/Fifty, $P_w(\langle P(H^n) \geq 2^{-n} \rangle) > 0$. By Proposition
 765 A.5, $P_w(H) < \frac{\ell(n)}{n}$. Since $\langle W, \mathcal{P} \rangle$ validates Simple Trust, $P_w(H^n \mid \langle P(H^n) \geq 2^{-n} \rangle) \geq 2^{-n}$.
 766 So, by Lemma A.6,

$$(7) \quad P^\ell(H^n \mid \langle P(H^n) \geq 2^{-n} \rangle) \geq 2^{-n}$$

Suppose $P^\ell(\langle P(H^n) \geq 2^{-n} \rangle) \geq 2^{-k}$. We have:

$$(8) \quad \begin{aligned} P^\ell(H^n \mid \langle P(H^n) \geq 2^{-n} \rangle) &= \frac{(\ell(n)/n)^n}{P^\ell(\langle P(H^n) \geq 2^{-n} \rangle)} \\ &\leq \frac{(\ell(n)/n)^n}{2^{-k}} \end{aligned}$$

So, from lines (7) and (8), it follows that:

$$2^k \left(\frac{\ell(n)}{n} \right)^n \geq 2^{-n}$$

which holds iff

$$\begin{aligned} \ell(n) &\geq \frac{n}{2^{1+\frac{k}{n}}} \\ &\geq \frac{n-k}{2} \end{aligned}$$

767 where the last line follows from Lemma A.7.

768

□

769 Having established a lower bound on $\ell(n)$, we now aim to establish an upper
 770 bound. The strategy is to consider a proposition true at just two worlds w and v
 771 (both IID), where $\#w = \frac{n}{2} - k$ and $\#v = \frac{n}{2} + k$. When k is sufficiently small, it will turn
 772 out that the proposition $\{w, v\}$ attains maximum probability amongst IID chances
 773 when $P = IID(1/2)$. This fact, which we establish in the next lemma, will then force
 774 IID chance functions at worlds with roughly $\frac{n}{2}$ occurrences of heads to assign heads
 775 probability $1/2$.

Lemma A.9. *Suppose n is even, $k \in \mathbb{N}$, and $k^2 \leq n/4$. Then the polynomial*

$$p^{\frac{n}{2}-k}(1-p)^{\frac{n}{2}+k} + p^{\frac{n}{2}+k}(1-p)^{\frac{n}{2}-k}$$

776 *achieves its maximum over the unit interval uniquely at $p = 1/2$.*

Proof. Without loss of generality, assume $p \in [0, 1/2]$. When $p = 1/2$, the polynomial
 evaluates to $2/2^n$, so we need to show

$$p^{\frac{n}{2}-k}(1-p)^{\frac{n}{2}+k} + p^{\frac{n}{2}+k}(1-p)^{\frac{n}{2}-k} \leq \frac{2}{2^n}$$

777 with equality iff $p = 1/2$. From simple algebra, we see this holds iff:

$$(9) \quad (2p)^{\frac{n}{2}-k}(2-2p)^{\frac{n}{2}+k} + (2p)^{\frac{n}{2}+k}(2-2p)^{\frac{n}{2}-k} \leq 2$$

778 Let $x = 1 - 2p$, so $x \in [0, 1]$. Line (9) holds just in case:

$$(1-x)^{\frac{n}{2}-k}(1+x)^{\frac{n}{2}+k} + (1-x)^{\frac{n}{2}+k}(1+x)^{\frac{n}{2}-k} \leq 2$$

779 This in turn holds iff:

$$(10) \quad (1-x)^{\frac{n}{2}-k}(1+x)^{\frac{n}{2}-k} \left[(1-x)^{2k} + (1+x)^{2k} \right] \leq 2$$

780 Further, the left-hand side of line (10) decreases with n . Since $k^2 \leq n/4$, we just need
781 to check that it holds for $n = 4k^2$.

782 The right- and left-hand sides are equal in line (10) when $x = 0$. The left-hand
783 side is differentiable, so to prove the theorem we just need to show the derivative
784 is negative.

785 Taking the derivative of the LHS of line (10) when $n = 4k^2$ and simplifying is
786 tedious, but we end up with:

$$-2k(1-x^2)^{2k^2-k-1} \left((1+x)^{2k}(2kx-1) + (1-x)^{2k}(2kx+1) \right)$$

787 Factoring out the $-2k(1-x^2)^{2k^2-k-1}$ out front, we see we need to verify that:

$$(11) \quad (1+x)^{2k}(2kx-1) + (1-x)^{2k}(2kx+1) > 0$$

788 for $k \geq 1$.

789 Using binomial expansion, we see verifying line (11) is equivalent to verifying:

$$(12) \quad \sum_{i=0}^{2k} \binom{2k}{i} \left[x^i(2kx-1) + (-x)^i(2kx+1) \right] > 0$$

790 The left-hand-side of line (12), in turn, simplifies to:

$$4kx^{2k+1} + 2 \sum_{i=0}^{k-1} \left[\binom{2k}{2i} 2k - \binom{2k}{2i+1} \right] x^{2i+1}$$

791 It is straightforward to check that $\binom{2k}{2i} 2k - \binom{2k}{2i+1} > 0$, which ensures the inequality
792 of line (12) holds, as desired. \square

793 We now can provide an upper bound on $\ell(n)$.

794 **Lemma A.10.** *Suppose $\langle W, \mathcal{P} \rangle$ is an n -flip model satisfying Simple Trust, Fifty/Fifty,
795 Symmetry, and Sufficiency with $n \geq 4$. Then $\ell(n) \leq \frac{n-\sqrt{n}}{2}$.*

796 *Proof.* Let $j \leq \frac{\sqrt{n}}{2}$ with $j \in \mathbb{N}$. By Sufficiency and Symmetry, there exist $w, v \in W$ such
797 that $\#w = \frac{n}{2} - j$ and $\#v = \frac{n}{2} + j$ and where P_w and P_v are both IID, and $P_v(T) = P_w(H)$.

798 Consider the proposition $X = \{w, v\}$. Let P_h be an IID chance function at a world
799 h with $\#h = n/2$. By Fifty/Fifty, $P_h(H) = 1/2$. Lemma A.9 entails that P_h assigns a
800 strictly higher probability to X (*viz*, 2^{-n+1}) than any other IID probability function
801 does.

802 *Claim:* $P_w(H) = 1/2$. For suppose not. Then $P_v(H) \neq 1/2$. In this case, $X \cap \langle P(X) \geq$
803 $2^{-n+1} \rangle = \emptyset$. So, since $P_h(\langle P(X) \geq 2^{-n+1} \rangle) > 0$, $P_h(X \mid \langle P(X) \geq 2^{-n+1} \rangle) = 0$, violating
804 Simple Trust.

805 So, if $j \leq \frac{\sqrt{n}}{2}$, then $\frac{n}{2} - j \leq \ell(n)$. Therefore, $\ell(n) \leq \frac{n-\sqrt{n}}{2}$ as desired. \square

806 **Theorem A.11.** *Suppose $\langle W, \mathcal{P} \rangle$ is an n -flip model that validates Simple Trust, Symmetry,
807 and Fifty/Fifty and $n \geq 6$. Suppose all functions in \mathcal{P} are IID. Then for all $w \in W$, if
808 $0 < \#w < n$, $P_w(H) = 1/2$.*

809 *Proof.* Suppose $\ell(n) \geq 1$, and let $P^\ell(H) = \frac{\ell(n)}{n}$ with P^ℓ an IID probability function
 810 defined over W . Let X be a random variable such that $X(w) = \#w$ for $w \in W$.
 811 $X \sim B(n, \ell(n)/n)$ according to P^ℓ . The mode of $B(n, \ell(n)/n) = \ell(n)/n$, which means $P^\ell(\langle X \geq$
 812 $\ell(n)/n \rangle) \geq \frac{1}{2}$. By Lemma A.8, $\ell(n) \geq \frac{n-1}{2n}$. Since $\ell(n)$ is an integer with $n \geq 6$, $\ell(n) = n/2$.
 813 But by Lemma A.10, $\ell(n) \leq \frac{\sqrt{n}}{2} - 1$. So, $n/2 \leq \frac{\sqrt{n}}{2} - 1$, which is impossible when
 814 $\ell(n) \geq 6$. So, $\ell(n) = 1$ for all $n \geq 6$. This completes the proof. \square

815 **Theorem A.12.** *Suppose $\langle W, \mathcal{P} \rangle$ is an n -flip model that validates Symmetry, Fifty/Fifty,
 816 and $n \geq 6$, and π is a regular probability function that totally trusts $\langle W, \mathcal{P} \rangle$. Suppose all
 817 functions in \mathcal{P} are IID. Then for all $w \in W$ if $0 < \#w < n$, $P_w(H) = 1/2$.*

818 *Proof.* This follows immediately from Theorems A.1 and A.11. \square

819 We will now see how we can relax the assumption that all chance functions are
 820 IID and still cause trouble for the Humeans.

821 **Lemma A.13.** *Suppose $\langle W, \mathcal{P} \rangle$ is an n -flip model satisfying Simple Trust, Fifty/Fifty,
 822 Monotonicity, Symmetry, and Sufficiency. Then $P^\ell(\langle P(H^n) \geq 2^{-n} \rangle) \leq 2^{-\sqrt{n}}$.*

823 *Proof.* Given the assumptions, we know from Lemma A.8, that if $P_w(\langle P(H^n) \geq$
 824 $2^{-n} \rangle) > 2^{-k}$, $\frac{n-k}{2} \leq \ell(n)$. From the assumptions and Lemma A.10, we know $\ell(n) \leq$
 825 $\frac{n-\sqrt{n}}{2}$. So, $k \geq \sqrt{n}$, meaning $P^\ell(\langle P(H^n) \geq 2^{-n} \rangle) \leq 2^{-\sqrt{n}}$. \square

826 Note that $\langle \text{IID}(P) \text{ and } P(H) \geq \frac{1}{2} \rangle \subseteq \langle P(H^n) \geq 2^{-n} \rangle$. So, what Lemma A.13 entails
 827 is the following: Let P_w be an IID chance function that assigns probability under
 828 $1/2$ to H , but such that if $P_v \in \mathcal{P}$ is IID and $P_v(H) < 1/2$, then $P_v(H) \leq P_w(H)$. It's
 829 easy to show, given the assumptions, that $P_w(\langle P(H^n) \geq 2^{-n} \rangle) < P^\ell(\langle P(H^n) \geq 2^{-n} \rangle)$.

830 Intuitively, at least when n is big, $P_w(H)$ should be *just* under $1/2$. After all,
 831 if just one more tail had been heads, then (if done in a way that maintained
 832 IID), the chance of heads would have been $1/2$. But Theorem A.13 entails that
 833 $P_w(\langle P(H^n) \geq 2^{-n} \rangle) \leq 2^{-\sqrt{n}}$, which is small. (E.g., when n is 10, this quantity is
 834 $P_w(\langle P(H^n) \geq 2^{-n} \rangle) < .12$. When $n = 100$, $P_w(\langle P(H^n) \geq 2^{-n} \rangle) < .001$.) This can only
 835 be the case if *either* $P_w(H)$ is extremely small, *or* very few worlds have IID chance
 836 functions. Indeed, as n grows, the proportion of worlds with approximately $n/2$
 837 heads tends toward 1 (where ‘‘approximately’’ here means within $x\%$ of $n/2$). So,
 838 either $P_w(H)$ must tend toward 0 or the percentage of worlds with IID chance
 839 functions must tend toward 0 very quickly. This is why, intuitively, when we add
 840 Boundedness, we end up with the Serious Triviality result in the main text.

841 **Theorem A.14.** *Let $\langle W_1, \mathcal{P}_1 \rangle, \langle W_2, \mathcal{P}_2 \rangle, \dots$ be a sequence of models with $|W_i| < |W_{i+1}|$.
 842 Assume each validates Simple Trust, Sufficiency, Fifty/Fifty, and Symmetry. Moreover,
 843 assume that Boundedness holds of the sequence. Then there exists an $N \in \mathbb{N}$ such that if
 844 $i \geq N$ and $P_w \in \mathcal{P}_i$ is IID, then $P_w = \text{IID}(1/2)$.*

845 *Proof.* Suppose $\ell(n) \geq 2$. Let P^ℓ be IID with $P^\ell(H) = \frac{\ell(n)}{n}$. Let $\text{IID}(W) := \{w \in W : P_w$
 846 $P_w \text{ is IID}\}$, and let $h(W) := \{w \in W : \ell(n) \leq \#w \leq n - \ell(n)\}$.

847 Note that, given Symmetry, if $w \in h(W) \cap \text{IID}(W)$, then $P_w(H) = 1/2$. So,

$$(13) \quad d \cdot P^\ell(h(W)) \leq P^\ell(h(W) \cap \text{IID}(W)) \leq 2^{-\sqrt{n}}$$

848 where the first inequality follows from Strong Sufficiency with threshold d , and the
 849 second from Lemma A.13.

850 We will now show that for large enough n , $d \cdot P^\ell(h(W)) > 2^{-\sqrt{n}}$, contradicting line
 851 (13). For fixed $\langle W, \mathcal{P} \rangle$, let $X(w) = \#w$. If $X \sim B(n, p)$, then X has increasing variance
 852 with p over $[0, 1/2]$. By Lemma A.10, $\ell(n) \leq \frac{n-\sqrt{n}}{2}$. So, the minimum possible value
 853 for $P^\ell(h(W))$ is achieved when $P^\ell(H) = \frac{n-\sqrt{n}}{2n}$.

854 So, assume $P^\ell(H) = \frac{n-\sqrt{n}}{2n}$. If $X \sim B\left(n, \frac{n-\sqrt{n}}{2n}\right)$, then $\sigma(X) = \frac{\sqrt{n-1}}{2}$, where $\sigma(X)$ rep-
 855 represents the standard deviation of X . By Chebyshev's Inequality, we then know
 856 $P^\ell\left(\frac{n-3\sqrt{n}}{2} \leq X \leq \frac{n+\sqrt{n}}{2}\right) > \frac{3}{4}$ (since the probability X is within two standard devia-
 857 tions must be at least $\frac{3}{4}$). But, the mode of X is $\ell(n)$, so $P^\ell(X < \ell(n)) < 1/2$. Therefore,
 858 $P^\ell\left(\ell(n) \leq X \leq \frac{n+\sqrt{n}}{2}\right) > 1/4$. Thus, $P^\ell(h(W) \cap \text{IID}(W)) > d/4$. For sufficiently large n ,
 859 $d/4 > 2^{-\sqrt{n}}$, which contradicts line (13). So, for large enough n , $\ell(n) = 1$.
 860 □

861 We now can state our final triviality result, referred to as Serious Triviality in the
 862 main text.

863 **Theorem A.15.** *Let $\langle W_1, \mathcal{P}_1 \rangle, \langle W_2, \mathcal{P}_2 \rangle, \dots$ be a sequence of models with $|W_i| < |W_{i+1}|$.
 864 Assume each validates Sufficiency, Fifty/Fifty, and Symmetry. Moreover, assume that Bound-
 865 edness holds of the sequence. Then there exists an $N \in \mathbb{N}$ such that if $i \geq N$ and some
 866 regular probability function totally trusts $\langle W_i, \mathcal{P}_i \rangle$, then for all $P_w \in \mathcal{P}_i$ such that $\text{IID}(P_w)$,
 867 we have $P_w = \text{IID}(1/2)$.*

868 *Proof.* This follows from Theorems A.1 and A.14. □

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